

A Novel Design of Reversible Logic based 1-Bit and 4-Bit ALU

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Abstract— in last few years, the concept called Reversible logic (RL) has received more attention because of its several abilities like minimizing the power dissipation. Generally, the more power dissipation will take place in low powered digital designs. The applicability of RL has extended in several areas like design of low power CMOS, DNA computing, advanced computing, quantum computing, processing of optical information and in nanotechnology. In case of traditional digital circuits designs the major amount of power loss take place during the logic operation by which some information bits will be erased. Thus, in order to protect this information bits from erasing; the logic gates are needed to be designed. Generally, in reversible computation the information bits will not be erased. This concept of reversible computation gives the hint for design of RL gates. The Central Processing Unit (or CPU) has got a fundamental building block called as Arithmetic logic Unit (or ALU). The use of reversible ALU in a computing system can offer higher power optimization. The realization of arithmetic operations can be done by adopting proper control logic over any one input variable parallel to adder. In this paper, 1 bit reversible and 4bit ALU is presented by using RL. The designed ALU's are analyzed over the FPGA SPARTAN6 device. The performance of the proposed mechanism is compared with other existing mechanism by considering the parameters like garbage outputs, propagation delay and quantum cost.

Keywords- Arithmetic Logic Unit, central processing unit (CPU), Reversible logic, FPGA.

I. INTRODUCTION

In design of digital circuits, the energy dissipation is always an biggest concern. One of the reasons connected to the power dissipation in digital circuit is technical non-ideality of the materials and switches. Generally, the logic operations which performed over the traditional digital circuit designs lead higher energy dissipation, information bits loss and heat loss. Thus in order to overcome these issues related to the digital designs, several fabrication, reverse computation and integration mechanism are exist. The use of reversible computation can provide the protection against the loss of information bits and also minimize the power dissipation. The logic of reversible computation has planted the development of reversible logic (RL). Thus, since last decade the reversible logic has given promise to resolve the power dissipation issues. The RL has applicability in various application sectors like design of low power CMOS, DNA computing, advanced computing, quantum computing, processing of optical

information and in nanotechnology. The non use of RL concept or use of irreversible hardware computation will lead in loss of energy because of information bits loss/erase. The researcher Launderer have suggested that the energy dissipation in irreversible bit operation will be $KT(\ln(2))$ Joules, where K is Boltzmann's constant and its value is $1.38065 \times 10^{-23} \text{K(joule/Kelvin-1)}$ and T is the temperature during the operation. Also its is says that the energy dissipation can be controlled or eliminated by making use of RL forward and backward operation, which means the reversible computation will produce inputs through outputs and at any point of computation it can stop and returns to history [1].

The analyst Bennett demonstrated that the dissipated power is straightly identified with the number of bits which were erased during process., furthermore that the computers can be reversible logically, lessened complexity and at helpful speed create précised calculations and to keep away from $kT(\ln 2)$ joules of energy dissipation in a circuit it ought to be produced using RLC. For this the circuit must be logically reversible. Another approach of configuration comes in the field of digital circuits outlining for restricting the power dispersal. The device designed by new approach is known as a Reversible Logic Device (RLD). An entryway designed utilizing RL is called Reversible Logic Gate (RLG). Inputs are dictated by the programmers for execution in a guideline set architecture. In view of this input number of ALU to have the capacity to produce assortment of logic outputs. Consequently in this sort of environment RLC must have both settled select input lines that get opcode signals controlled by programmed and output lines where the logical output result is delivered [2].

RLG is an *I –input - O –Output* logic function. A RLG is also defined as a bijective Boolean function from *I* to *O* values. Let the input vector be P_v –input vector , Q_v –output vector and they are defined as follows,

$$P_v = (P_i, P_{i+1}, \dots, P_{I-1})$$

$$Q_v = (Q_i, Q_{i+1}, \dots, Q_{I-1}, Q_I)$$

This will satisfy relation

$$P_v = Q_v$$

In the design of RLG the below points are necessary.

- Low Garbage outputs
- Low delay.
- There should not be feedbacks or loops
- There should not be Fan out
- Low quantum cost.

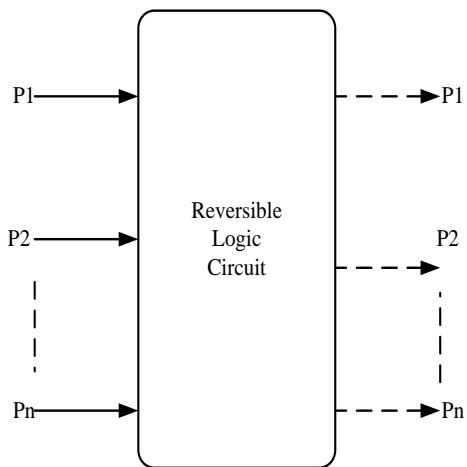


Fig.1. Reversible Logic gate (RLG) with $I - input$ and $O - Output$

The paper is organized as follows section II highlights prior research work pertaining to Reversible ALU. Section III describes problems explored after reviewing the literatures. Discussion of the proposed system using novel approach is given in section IV. The Section V will demonstrate the outcomes of the proposed study followed by Conclusion of paper in section VI.

II. RELATED WORK

Our earlier survey has as of now talked about different existing framework to Design of Reversible Logic ALU utilizing RLG [3] with an unmistakable depiction of the exploration crevice.

This segment, we audit some more significant research work towards RL ALU. In Zhijin Guan et al. [4] a plan building the ALU in view of RLG as rationale segments is proposed. The exhibited reversible ALU lessens the information bits' utilization and misfortune by reusing the rationale information bits logically and understands the objective of bringing down power utilization.

In Gopal et al. [5] paper, two sorts of reversible ALU outlines are proposed and checked utilizing Altera Quartus II programming. In the proposed outlines, eight number juggling and four logical operations are performed. In the proposed plan 1, Peres Full Adder Gate (PFAG) is utilized as a part of reversible ALU outline and HNG door is utilized as a viper rationale circuit in the proposed ALU outline 2. Both proposed plans are dissected and looked at as far as number of entryways check, junk output, quantum cost and spread postponement. The reproduction results demonstrate that the

proposed reversible ALU outline 2 out forms the proposed reversible ALU plan 1 and ordinary ALU plan.

In this survey of Praveen Kumar [6] have explored that the Programmable RL is developing as a forthcoming RL for usage in quantum computing and cutting edge nanotechnology with insignificant effect on circuit warm era.

Premanand and Ravindranath [7] portrayed the plan of programmable RLG structures, focused for the ALU execution using an effective reversible ALU. Utilizing reversible rationale doors rather than customary rationale AND/OR entryways, a reversible ALU whose capacity is the same as conventional ALU is developed. Contrasting and the input bits and the lost of bits of the conventional ALU, the reversible ALU essentially decrease the utilization and loss of information bits. Programmable RLG are acknowledged in Verilog HDL, reproduced and orchestrated utilizing Cadence NCSIM and RTL compiler.

Hence, it can be observed that different techniques being introduced in recent time for evolving up with new strategies of reversible logic ALU. All the techniques discussed have significant beneficial point of learn and adopt while associated with limitations and constraints too. The problems pertaining to existing studies are discussed briefly in next section.

III. PROBLEM DESCRIPTION

This section presents the problems that have been identified after reviewing the existing techniques on RL in ALU. It has been found that focuses of existing techniques on RL in ALU are quite less. Design of a control unit for any figuring unit is the hardest part and includes more basic limitations. Power consumption is a critical issue in VLSI designs. The progression in VLSI designs and especially compact gadget advances and progressively high computation necessities, prompt to the design of speedier, small and more perplexing electronic Systems. The coming of multi-giga-hertz processors, top of the line electronic devices carry with them an expansion in framework complexity, packets with high density and a issue on power consumption. Hence, the next section discusses about proposed system to overcome this issues.

IV. PROPOSED FRAMEWORK

The design of 1-bit ALU and 4-bit ALU utilizing reversible logical gate is simulated using Verilog. The realized ALU is examined on FPGA device. This work is compared in terms of quantum cost, the propagation delay and the garbage outputs. Here, we designed a 4-bit reversible logic ALU and also the implementation on proposed 1-bit reversible ALU architecture. The main part of the ALU is the parallel adder. It is designed with a many number of full adders and also the projected 1-bit as well as the 4-bit reversible ALU and it also uses DKFG gate is used as a full adder by calculating the inputs to the parallel adder. It is probably to obtain the different type of ALU operations. The carry input C_{in} is applied to the full adder circuit in the lower bit position and it

exits the Cout as of the adder. The full adder is used to obtain the sum of gets and carry output.

The functional expression of an ALU is given by

$$F_i = X_i \otimes^i \otimes \text{ and } C_{i+1} = X_i Y_i + Y_i Z_i + X_i Z_i$$

Where X_i, Y_i and Z_i are given by, $X_i = A_i + S_2 S_0 (S_1 B_i) \otimes$

$$Y_i = S_0 B_i + S_1 B_i$$

$$Z_i = S_2 C_i$$

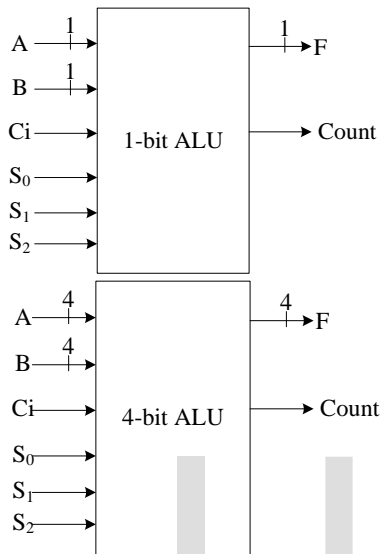


Fig.2. Top Module

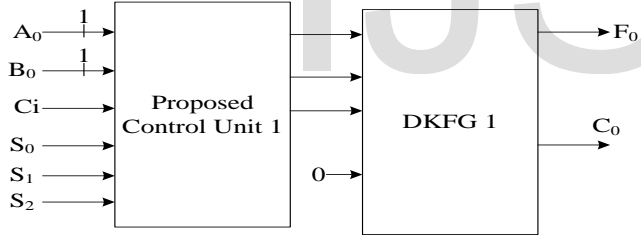


Fig.3. Top Module

Figure 2. Shows the block diagram of proposed 1-bit as well as 4-bit ALU model. It contains the three select inputs such as S1, S2 and S3. Where, the Ai, Bi and Ci are the 3 different inputs. Zi, Yi and Xi are the output obtained from proposed control unit. These signals are given to the DKFG full adder circuit to perform the summing operations. Fi and Cout are the two output obtained from full adder circuit. Table 1 shows the truth table of the proposed 1-bit ALU model.

Form this section we can conclude that, the delay, garbage value and the quantum cost obtained in 4-bit ALU is more than 1-bit ALU. The projected ALUs has mainly two most significant advantages. Firstly, it obtains less propagation delay is compared with existing 1-bit architecture. As a result, the novel and newly developed ALU has a best delay for n-bit ALUs operations. Secondly, the work of ALU generates more number of arithmetical as well as logical operations. The projected ALU is well designed and it is flexible that, any type

of alterations actions required for simulation in instruction data set architecture will be made easy.

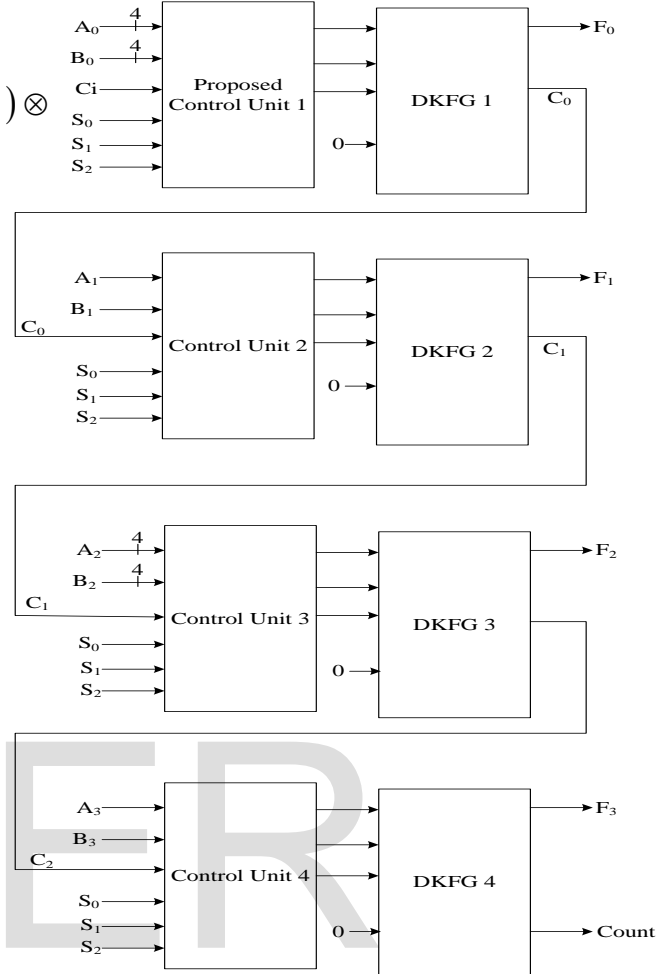


Fig.4. Internal Architecture of 4-bit reversible ALU

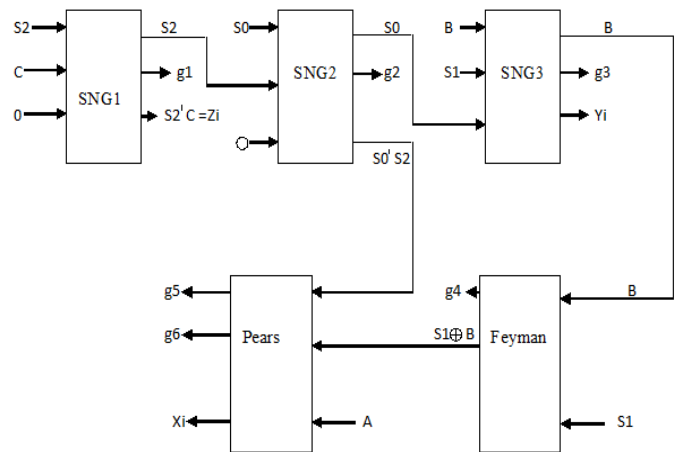


Fig.5. Proposed Control Unit for ALU

Table.1. The truth table of the proposed method

S2	S1	S0	Cin	Output	Function
0	0	0	0	F=A	Transfer
0	0	0	1	F=A+1	Increment
0	0	1	0	F=A+B	Addition
0	0	1	1	F=A+B+1	Addition with Carry
0	1	0	0	F=A'+B	Subtraction with Borrow
0	1	0	1	F=A'+B+1	Subtraction
0	1	1	0	F=A-1	Decrement
0	1	1	1	F=A	Transfer
1	0	1	X	F=AUB	OR
1	1	0	X	F=⊗ AB	XOR
1	1	0	X	F=A ∧ B	AND
1	1	1	1	F=A	Complement

V. RESULTS AND DISCUSSION

A. Results of 1-bit ALU

The simulated 1-bit as well as 4-bit ALU framework illustrated and demonstrated in the previous part is amalgamated and simulated on Xilinx. The 1-bit ALU top module is shown in fig.6. It having three inputs such as a, b, cin. Then the select line input is take as s1, s2 and s3 respectively. The output of this module is cout and function.

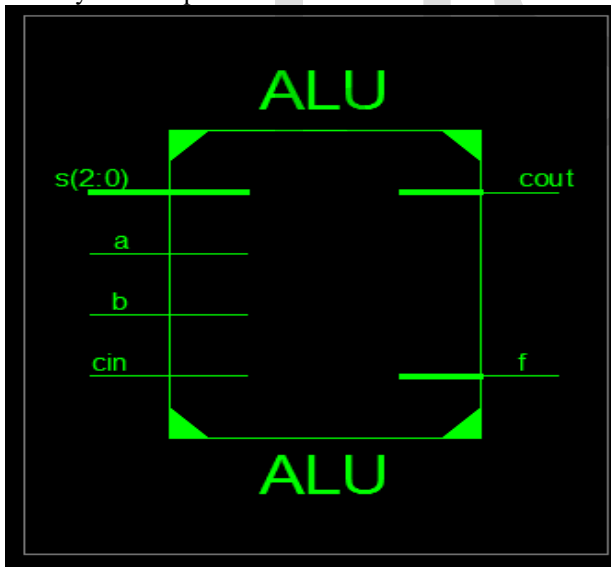


Fig.6. ALU_1-bit_top_module

The RTL diagram proposed 1-bit ALU as show in figure. 6. It is the inbuilt architect ere of the 1-bit ALU module. Similarly, the technological schematic 1 bit ALU diagram shown in fig.7. This shows the design and development of the 1-bit ALU using Verilog software.

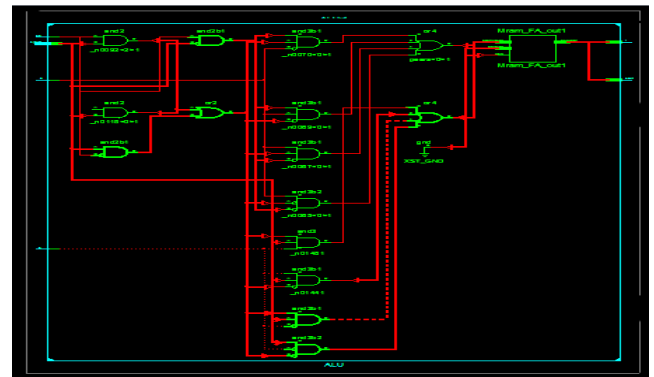


Fig.7. ALU_1-bit_rtl

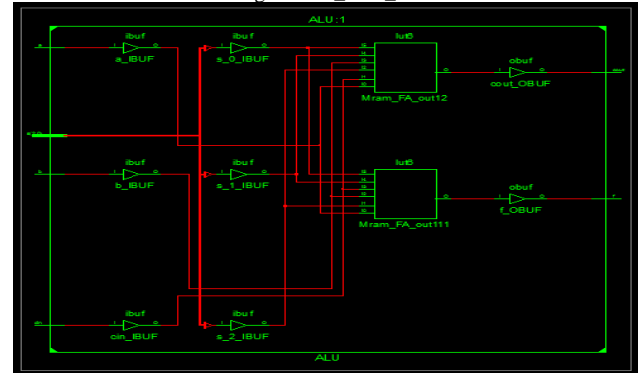


Fig.8. ALU_1-bit_technology_schematic

The simulated results also shown in fig.9. This process uses the following elements for the simulation purposes. This proposed method uses the Artix-7 as a FPGA circuit, it also uses a device like XC7A100T and Xc7a100t-3csf324 device, the package is CSG324 and the speed of the device is -3. It calculate the number of gate count, number of garbage and number of constant input is given to this simulation work.

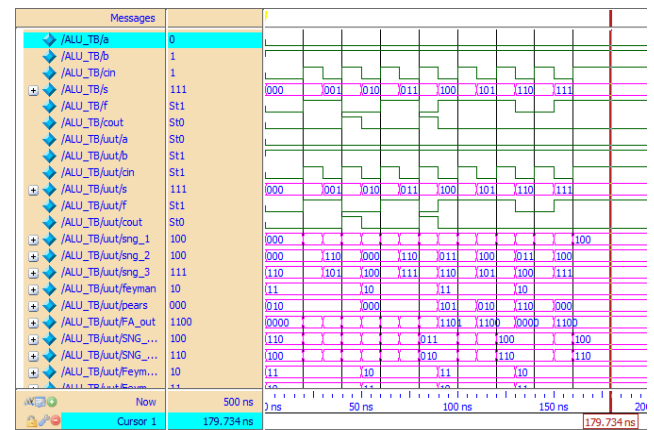


Fig.9. ALU_1-bit_Simulation

The summery of the device utilization values are estimated as shown in table. 1. It gives the usage of the devices for simulation work. The number of slice of LUTs is used as 2 and the available is 63,400. Similarly, the number of fully utilized LUT-FF pairs is zero and the available is 2. And finally, the number of bonded IOBs is used 8 and the available is 210. So, the utilization bonded IOBs are 3%, remaining are

0%. The target device is xc7a100t-3csg324, the version of product is ISE 14.7 and it uses the Xilinx by default tool.

Table.2. Device utilization summary

Device Utilization Summary (Estimated Values)			
Logic of Slice LUTs	Used	Available	Utilization
No. of Slice LUTs	2	63400	0%
No. of Fully used LUT-FF Pairs	0	2	0%
No. of Bonded IOBs	8	210	3%

From this 1-bit ALU design we can observed that, the maximum combination path delay is 1.075ns.

B. Results of 4-bit ALU

The simulated 4-bit ALU framework illustrated and demonstrated in the previous part is amalgamated and simulated on Xilinx. The 4-bit ALU top module is shown in fig.10. It having three inputs such as a, b, cin. Then the select line input is taking as s1, s2 and s3 respectively. The output of this module is cout and function.

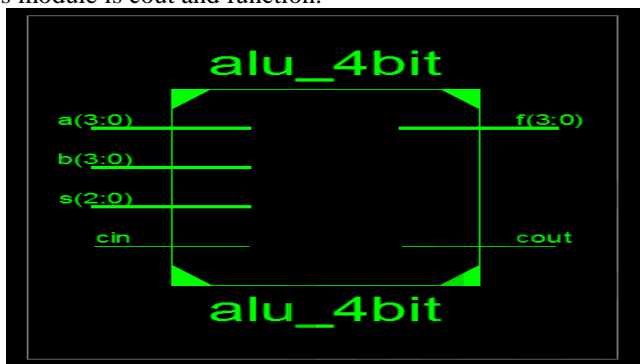


Fig.10. ALU_4-bit_top_module

The RTL diagram proposed 4-bit ALU as show in figure. 11. It is the inbuilt architect ere of the 4-bit ALU module. Similarly, the technological schematic diagram of the 4-bit ALU shown in fig.12. This shows the design and development of the 4-bit ALU using Verilog software.

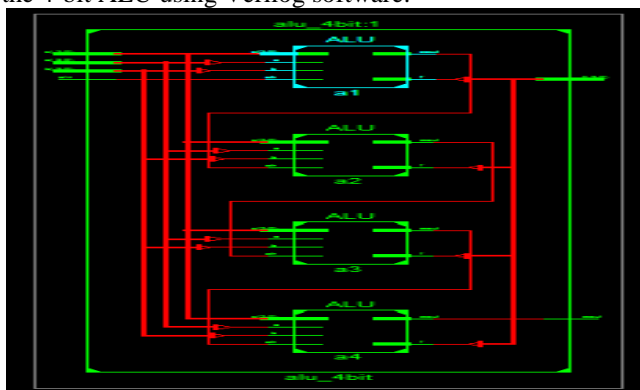


Fig.11. ALU_4-bit_rtl

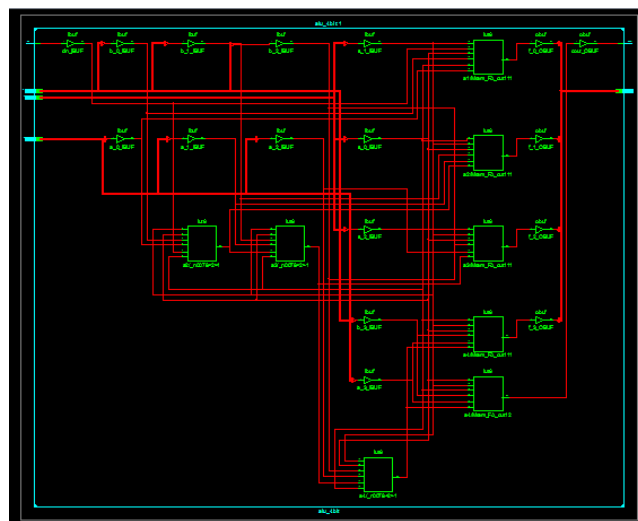


Fig.12. ALU_4-bit_technology_schematic

The simulated results also shown in fig.13. This process uses the following elements for the simulation purposes. This proposed method uses the Artix-7 as a FPGA circuit, it also uses a device like XC7A100T and Xc7a100t-3csg324 device, the package is CSG324 and the speed of the device is -3. It calculate the number of gate count, number of garbage and number of constant input is given to this simulation work.

The summary of the device utilization values are estimated as shown in table. 1. It gives the usage of the devices for simulation work. The number of slice of LUTs is used as 2 and the available is 63,400. Similarly, the number of fully utilized LUT-FF pairs is zero and the available is 2. And finally, the number of bonded IOBs is used 8 and the available is 210. So, the utilization bonded IOBs is 3%, remaining are 0%. The target device is xc7a100t-3csg324, the version of product is ISE 14.7 and it uses the Xilinx by default tool.

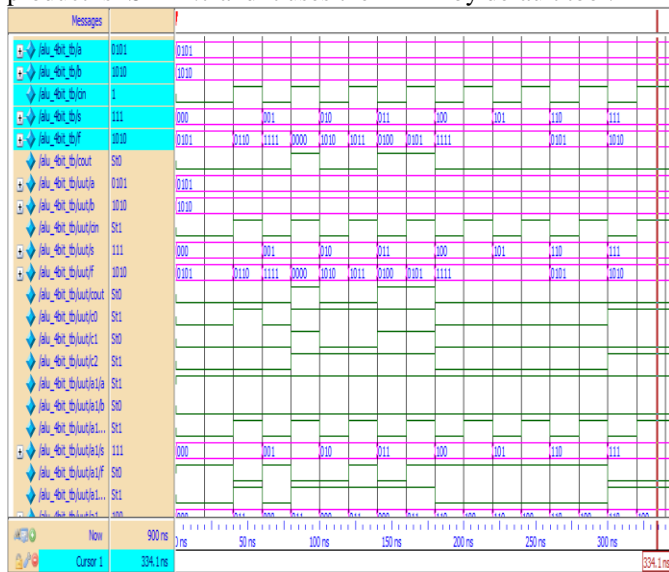


Fig.13. ALU_4-bit_Simulation

The summary of the device utilization values are estimated as shown in table. 3. It gives the usage of the devices for

simulation work. The number of slice of LUTs is used as 8 and the available is 63,400. Similarly, the number of fully utilized LUT-FF pairs is zero and the available is 8. And finally, the number of bonded IOBs is used 17 and the available is 210. So, the utilization bonded IOBs are 8%, remaining is 0%. The target device is xc7a100t-3csg324, the version of product is ISE 14.7 and it uses the Xilinx by default tool.

Table.3. Device utilization summary

Device Utilization Summary (Estimated Values)			
Logic of Slice LUTs	Used	Available	Utilization
No. of Slice LUTs	8	63400	0%
No. of Fully used LUT-FF Pairs	0	8	0%
No. of Bonded IOBs	17	210	8%

From this 4-bit ALU design we can observed that, the maximum combination path delay is 3.466 ns.

Table.4. Evaluation among the projected reversible 1-bit and 4-bit ALU Schemes with the present 1-bit ALU proposal

Parameter which has to be compared	Proposed 4-bit reversible ALU	Proposed 1-bit reversible ALU	Existing 1-bit reversible ALU
Delay	7.885ns	6.334ns	16ns
Quantum Cost	88	22	24
Garbage Output	10	6	10

Table 4 gives the comparison between the projected 1-bit and 4-bit ALU using reversible logic circuit as well as the existing ALU results.

VI. CONCLUSION

The programmable RLG have been examined in reversible ALU. These new ALU designs are profitable to existing work in usage that support low delay and high logical output, which is attractive for acknowledgment of a reversible CPU. The proposed plans can be incorporated in the outline of an n-bit reversible ALU. In this paper examination and investigation done on gates on FPGA SPARTAN6 device. With the above investigation it can likewise be determined that proposed 1-Bit ALU has less postponement than the current 1-Bit ALU. From this we can observed that our proposed method gives better and good performance results. The propagation delay, the quantum cost and finally the garbage outputs are deliberated as the performance parameters. For conventional RLC there exists much research, even entire books, performed to the outline and simulation of computer arithmetic. This is unquestionably not the situation for RL. The requirement that the circuits must be garbage-free is the thing that makes it a fascinating exploration issue, however most proposed designs still actualize the conventional algorithms with garbage. In this utilized the RLG, however as their sole objective is to decrease logic size or

number of garbage-bits for a particular fixed size circuit; next better results are obtained from this approach.

VII. REFERENCES

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